

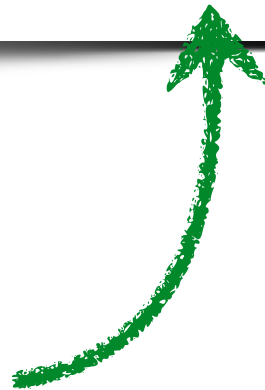
# Bounded Refinement Types

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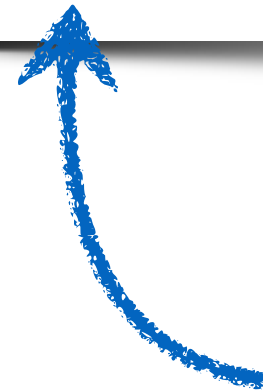
# Refinement Types

$\{v:\text{Int} \mid v > 0\}$

Basic Type



Predicate



## **LiquidHaskell:**

Refinement types to express specifications

## **Choice:**

Refinements drawn from decidable logic

## **Question:**

How to enhance expressiveness?

## **Question:**

How to enhance expressiveness?

**Can we specify function composition?**

# Can we specify function composition?

```
incr      :: x:Int -> {v | v = x+1}  
incr x = x+1
```

```
incr2 :: x:Int -> {v | v = x+2}  
incr2 = compose incr incr
```

# Can we specify function composition?

Type is too specific

```
compose ::  
    (y:b -> {z:c | z = y+1})  
  -> (x:a -> {z:b | y = x+1})  
  -> (x:a -> {z:c | z = x+2})  
compose f g x = f (g x)
```

```
incr2 :: x:Int -> {v | v = x+2}  
incr2 = compose incr incr
```

OK

# Can we specify function composition?

Type is too specific

compose ::

(y:b -> {z:c | z = y+1})

-> (x:a -> {z:b | y = x+1})

-> (x:a -> {z:c | z = x+2})

compose f g x = f (g x)

incr2 :: x:Int -> {v | v = x+2}

incr2 = compose decr incr

Fail

# Can we specify function composition?

Type is too specific

```
compose ::  
    (y:b -> {z:c | z = y+1})  
  -> (x:a -> {z:b | y = x+1})  
  -> (x:a -> {z:c | z = x+2})  
compose f g x = f (g x)
```

```
incr2 :: x:Int -> {v | v = x+2}  
incr2 = compose incr incr
```



# Can we specify function composition?

$p, q \mid \rightarrow \lambda x. z \rightarrow z = x + 1$   
 $r \mid \rightarrow \lambda x. z \rightarrow z = x + 2$

**compose** ::

$(y:b \rightarrow \{z:c \mid p \ y \ z\})$   
 $\rightarrow (x:a \rightarrow \{y:b \mid q \ x \ y\})$   
 $\rightarrow (x:a \rightarrow \{z:c \mid r \ x \ z\})$

**compose**  $f \ g \ x = f \ (g \ x)$

**incr2** ::  $x:\text{Int} \rightarrow \{v \mid v = x + 2\}$

**incr2** = **compose** **incr** **incr**

# Can we specify function composition?

Type is wrong

```
compose ::  
    (y:b -> {z:c | p y z})  
  -> (x:a -> {y:b | q x y})  
  -> (x:a -> {z:c | r x z})  
compose f g x = f (g x)
```

# Can we specify function composition?

Type is wrong

compose ::

(y:b -> {z:c | p y z})

-> (x:a -> {y:b | q x y})

-> (x:a -> {z:c | r x z})

compose f g x = let y = g x in f y

y : {y:b | q x y} ⊢ {z:c | p y z} <: {z:c | r x z}

# Can we specify function composition?

Type is wrong

compose ::

( $y:b \rightarrow \{z:c \mid p \ y \ z\}$ )

$\rightarrow (x:a \rightarrow \{y:b \mid q \ x \ y\})$

$\rightarrow (x:a \rightarrow \{z:c \mid r \ x \ z\})$

compose f g x = let y = g x in f y

$y:\{y:b \mid q \ x \ y\} \vdash \{z:c \mid p \ y \ z\} <: \{z:c \mid r \ x \ z\}$

# Can we specify function composition?

Type is wrong

compose ::

(y:b -> {z:c | p y z})

-> (x:a -> {y:b | q x y})

-> (x:a -> {z:c | r x z})

compose f g x = let y = g x in f y

y:{y:b | q x y} ⊢ {z:c | p y z} <: {z:c | r x z}

# Can we specify function composition?

Type is wrong

compose ::

=> (y:b-> {z:c | p y z})

-> (x:a-> {y:b | q x y})

-> (x:a-> {z:c | r x z})

compose f g x = let y = g x in f y

bound Chain p q r = \x y z ->

q x y => p y z => r x z

# Can we specify function composition?

`compose` :: (Chain p q r)  
=> (y:b-> {z:c | p y z})  
-> (x:a-> {y:b | q x y})  
-> (x:a-> {z:c | r x z}) **OK**

`compose` f g x = let y = g x in f y

bound Chain p q r = \x y z ->  
q x y => p y z => r x z

# Can we specify function composition?

$$\begin{array}{lcl} p, q & |-> & \backslash x \ z \rightarrow z = x + 1 \\ r & |-> & \backslash x \ z \rightarrow z = x + 2 \end{array}$$

`incr2` :: `x:Int` -> `{v | v = x+2}`  
`incr2` = `compose incr incr`

`bound Chain p q r = \x y z ->`  
`q x y => p y z => r x z`



# Can we specify function composition?

$$\begin{array}{lcl} p, q & |-> & \backslash x \ z -> z = x + 1 \\ r & |-> & \backslash x \ z -> z = x + 2 \end{array}$$

`incr2` :: `x:Int` -> {`v` | `v = x+2`}

`incr2` = `compose incr incr`

**OK**

`bound Chain` = `\x y z ->`  
`y=x+1 => z=y+1 => z=x+2`

**Valid**

# **Can we specify function composition?**

Bounds let us specify function composition

## **Do bounds add complexity?**

No. Bounds are desugared to unbounded types

# Bounds are desugared to unbounded types

```
compose :: (Chain p q r)
        => (y:b-> {z:c | p y z})
        -> (x:a-> {y:b | q x y})
        -> (x:a-> {z:c | r x z})
```

```
compose f g x =
  let y = g x in
  let z = f y in z
```

```
bound Chain p q r = \x y z ->
  q x y => p y z => r x z
```

# Bounds are desugared to unbounded types

```
compose :: $chain:(tchain p q r)
      -> (y:b-> {z:c | p y z})
      -> (x:a-> {y:b | q x y})
      -> (x:a-> {z:c | r x z})
```

```
compose $chain f g x =
```

```
  let y = g x in
```

```
  let z = f y in z
```

```
type tchain p q r = x:a -> y:b -> z:c ->
  {v | q x y => p y z => r x z}
```

# Bounds are desugared to unbounded types

```
compose :: $chain:(tchain p q r)
        -> (y:b-> {z:c | p y z})
        -> (x:a-> {y:b | q x y})
        -> (x:a-> {z:c | r x z})
```

```
compose $chain f g x =
```

```
  let y = g x in
```

```
  let z = f y in
```

```
  let _ = $chain x y z in z
```

OK

```
type tchain p q r = x:a -> y:b -> z:c ->
  {v | q x y => p y z => r x z}
```

# Bounds are desugared to unbounded types

$p, q$	$\vdash$	$\lambda x. z \rightarrow z = x + 1$
$r$	$\vdash$	$\lambda x. z \rightarrow z = x + 2$

```
incr2 :: x:Int -> {v | v = x+2}
incr2 = compose $chain incr incr
where $chain :: (tchain p q r)
      $chain = ???
```

```
type tchain p q r = x:a -> y:b -> z:c ->
  {v | q x y => p y z => r x z}
```

# Bounds are desugared to unbounded types

$p, q$	$\vdash$	$\lambda x. z \rightarrow z = x + 1$
$r$	$\vdash$	$\lambda x. z \rightarrow z = x + 2$

```
incr2 :: x:Int -> {v | v = x+2}
incr2 = compose $chain incr incr
where $chain :: tchain
      $chain = ???
```

```
type tchain      = x:a -> y:b -> z:c ->
  {v | y=x+1 => z=y+1 => z=x+2}
```

# Bounds are desugared to unbounded types

$p, q \mid -> \backslash x \ z -> z = x + 1$   
 $r \mid -> \backslash x \ z -> z = x + 2$

```
incr2 :: x:Int -> {v | v = x+2}
incr2 = compose $chain incr incr
where $chain :: tchain
      $chain = ???
```

```
type tchain      = x:a -> y:b -> z:c ->
  {v | true}
```



# Bounds are desugared to unbounded types

$p, q \mid -> \backslash x \ z -> z = x + 1$   
 $r \mid -> \backslash x \ z -> z = x + 2$

```
incr2 :: x:Int -> {v | v = x+2}
incr2 = compose $chain incr incr
where $chain :: tchain
      $chain = ???
```

```
type tchain      = x:a -> y:b -> z:c ->
  {v | true}
```

**Bounds enhance expressiveness**

**Do bounds add complexity?**

No. Bounds are desugared to unbounded types

**Are bounds useful?**

Function Composition

List Filtering    and    List Folding

# List Filtering

```
filter :: (Witness p w)
      => (x:a -> {v:Bool | w x v})
      -> [a] -> [{v:a | p v}]
filter q (x:xs)
  | q x      = x : filter q xs
  | otherwise = filter q xs
filter _ []  = []
```

```
bound Witness p w = \x b ->
  b => w x b => p x
```

# List Filtering

`isPos :: x:Int -> {v | v <=> 0 < x}`

`ex :: x:[Int] -> [{v | 0 < v}]`

`ex = filter isPos`

`p -> \x -> 0 < x`

`w -> \x b -> b <=> 0 < x`

`bound Witness p w = \x b ->`

`b => w x b => p x`

# List Filtering

`isPos :: x:Int -> {v | v <=> 0 < x}`

`ex :: x:[Int] -> [{v | 0 < v}]`

`ex = filter isPos`

$p \rightarrow \lambda x \rightarrow 0 < x$

$w \rightarrow \lambda x \ b \rightarrow b \iff 0 < x$

`bound Witness = \x b ->`

`b ==> (b <=> 0 < x ) ==> 0 < x`

**Valid**

**Bounds enhance expressiveness**

**Do bounds add complexity?**

No. Bounds are desugared to unbounded types

**Are bounds useful?**

Function Composition

List Filtering and List Folding

# List Folding

```
foldr :: (Inductive inv step)
=> (x:a -> acc:b -> {v:b | step x acc v})
-> {v:b | inv [] v} -> xs:[a]
-> {v:b | inv xs v}

foldr f b (x:xs) = f x (foldr op b xs)
foldr f b []     = b
```

```
bound Inductive inv step = \x xs b b' ->
  inv xs b => step x b b'
  => inv (x:xs) b'
```

# List Folding

`incr :: x:Int -> {v|v = x+1 }`

`ex :: xs:[a] -> {v:Int | v = len xs}`

`ex = foldr (\x -> incr) 0`

`inv -> \xs b -> b = len xs`

`step -> \x b b' -> b' = b + 1`

`bound Inductive inv step = \x xs b b' ->`

`inv xs b => step x b b'`

`=> inv (x:xs) b'`



# List Folding

`incr :: x:Int -> {v|v = x+1 }`

`ex :: xs:[a] -> {v:Int | v = len xs}`

`ex = foldr (\x -> incr) 0`

`inv -> \xs b -> b = len xs`

`step -> \x b b' -> b' = b + 1`

bound **Inductive** = `\x xs b b' ->`

`b = len xs => b' = b + 1`

`=> b' = len (x:xs)`

**Valid**

**Bounds enhance expressiveness**

**Do bounds add complexity?**

No. Bounds are desugared to unbounded types

**Are bounds useful?**

Function Composition

List Filtering and List Folding

Floyd-Hoare Logic in the State monad

Relational DataBases

***Thank you!***

**END**