SAT Solvers

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Decision Procedures

We will look very closely at the following

- 1. Propositional Logic
- 2. Theory of Equality
- 3. Theory of Uninterpreted Functions
- 4. Theory of Difference-Bounded Arithmetic

Decision Problem: Satisfaction

- Does eval s p return True for some assignment s ?
- "Can we assign the variables to make the formula true" ?

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Why?

- Representative
- Have "efficient" algorithms

Decision Procedures

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Plan

- First in isolation
- Then in combination
- Very slick SW-Eng, based on logic

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Popularly called SAT Solvers

Basics

- Propositional Logic 101
- Conjunctive Normal Form
- Resolution

Algorithms

- Resolution
- Backtracking Search
- Boolean Constraint Propagation
- Conflict Driven Learning & Backjumping

Basics

Propositional Logic 101

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Propositional Logic 101

Propositional Variables

data PVar

Propositional Formulas

```
data Formula = Prop PVar
| Not Formula
| Formula 'And' Formula
| Formula 'Or' Formula
```

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Conjunctive Normal Form

Restricted representation of Formula

Literals: Variables or Negated Variables

data Literal = Pos PVar | Neg PVar

Clauses: Disjunctions (Or) of Literals

data Clauses = [Literal]

CNF Formulas: Conjunctions (And) of Clauses

data CnfFormula = [Clauses]

Conjunctive Normal Form: Example

Consider a Formula $(x_1 \lor x_2) \land (\neg x_1 \lor x_3) \land \neg x_3$

Represented as a Formula

(Prop 1 'Or' Prop 2) 'And' (Not (Prop 1) 'Or' Prop 3) 'And' (Not (Prop 3))

Represented as a CnfFormula

```
[ [Pos 1 , Pos 2]
, [Neg 1 , Pos 3]
, [Neg 3 ]]
```

Conjunctive Normal Form Conversion

Theorem There is a *poly-time* function

```
toCNF :: Formula -> CnfFormula
toCNF = error "Exercise For The Reader"
```

Such that any f is satisfiable *iff* (toCNF f) is satisfiable.

- toCNF adds new variables for sub-formulas
- otherwise, an exponential blowup in CnfFormula size

Conjunctive Normal Form Conversion

Theorem There is a *poly-time* function

toCNF :: Formula -> CnfFormula
toCNF = error "Exercise For The Reader"

Such that any f is satisfiable *iff* (toCNF f) is satisfiable.

Henceforth Only consider formulas in Conjunctive Normal Form Formulas

Basics

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Properties of CNF

Pure Variable

▶ One which appears only +ve or -ve in a CnfFormula

Empty Clause

▶ If a CnfFormula has some Clause without Literals

Then the CnfFormula is UNSAT

Trivial Formula

- If a CnfFormula has no Clause
- Or every variable is pure
- Then the CnfFormula is SAT

Determine satisfaction by reducing CnfFormula to one of

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- Empty Clause (ie UNSAT), or
- Trivial Formula (ie SAT).

Reducing Formulas By Resolution

("Reduce" is, perhaps, not the best word...)

Resolution: For any A, B and variable x, the formula

 $(A \lor x) \land (B \lor \neg x)$

is equivalent to the formula

 $(A \lor B)$

The variable x is called a **pivot** variable

General Resolution

Resolution: For any A_i, B_j and variable x, the formula

$$\bigwedge_i (A_i \lor x) \land \bigwedge_j (B_j \lor \neg x)$$

is equivalent to the formula

$$\bigwedge_{i,j} (A_i \vee B_j)$$

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Pivot variable x is eliminated by resolution

Davis-Putnam Algorithm: Example 1

Input Formula

$$\blacktriangleright (x_1 \lor x_2 \lor x_3) \land (x_2 \lor \neg x_3 \lor x_5) \land (\neg x_2 \lor x_4))$$

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Pivot on x_2

$$\blacktriangleright (x_1 \lor x_3 \lor x_4) \land (\neg x_3 \lor x_5 \lor x_4)$$

Pivot on x₃

$$\blacktriangleright (x_1 \lor x_4 \lor x_5)$$

All variables are *pure* ... hence, **SAT**

Davis-Putnam Algorithm: Example 2

Input Formula

$$\blacktriangleright (x_1 \lor x_2) \land (x_1 \lor \neg x_2) \land (\neg x_1 \lor x_3) \land (\neg x_1 \lor \neg x_3)$$

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Pivot on x_2

$$\blacktriangleright (x_1) \land (\neg x_1 \lor x_3) \land (\neg x_1 \lor \neg x_3)$$

Pivot on x_3

$$\blacktriangleright (x_1) \land (\neg x_1)$$

Pivot on x_1

Empty clause ... hence, **UNSAT**

Davis-Putnam Algorithm

Algorithm

- 1. Select **pivot** and perform **resolution**
- 2. Repeat until SAT or UNSAT

Issues?

Space blowup (formula size blows up on resolution)

Basics

- Propositional Logic 101
- Conjunctive Normal Form

Algorithms

- Resolution
- Backtracking Search
- Boolean Constraint Propagation
- Conflict Driven Learning & Backjumping

Decision Tree: Describes Space of All Assignments

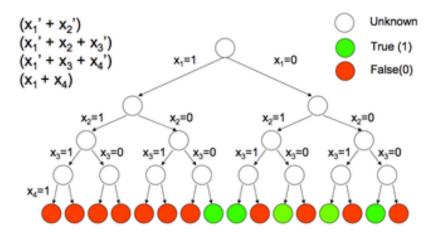


Figure: SAT Decision Tree (Courtesy: Lintao Zhang)

Decision Tree: SAT via Depth First Search

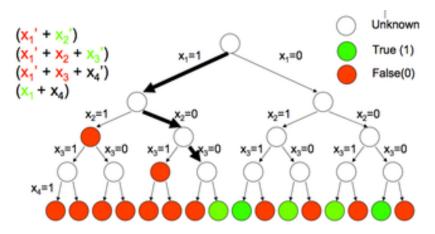


Figure: DFS On Decision Tree (Courtesy: Lintao Zhang)

Don't build whole tree, but lazily search solutions

- Choose a variable x, set to True
- Remove constraints where x appears
- Recurse on remaining constraints
- Backtrack if a contradiction is found

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Backtracking Search (1/21)

 $\begin{array}{l} (a'+b+c)\\ (a+c+d)\\ (a+c+d')\\ (a+c'+d)\\ (a+c'+d')\\ (b'+c'+d)\\ (a'+b+c')\\ (a'+b+c')\\ (a'+b'+c) \end{array}$

Figure: Basic DLL (Courtesy: Lintao Zhang)

Backtracking Search (2/21)

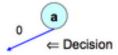


(a' + b + c) (a + c + d) (a + c + d') (a + c' + d) (a + c' + d') (b' + c' + d) (a' + b + c') (a' + b' + c)

Figure: Basic DLL (Courtesy: Lintao Zhang)

Backtracking Search (3/21)

| (a' + b + c) |
|---------------|
| (a + c + d) |
| (a + c + d') |
| (a + c' + d) |
| (a + c' + d') |
| (b' + c' + d) |
| (a' + b + c') |
| (a' + b' + c) |

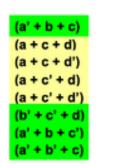


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Figure: Basic DLL (Courtesy: Lintao Zhang)

Backtracking Search (4/21)



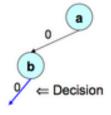


Figure: Basic DLL (Courtesy: Lintao Zhang)

Backtracking Search (5/21)

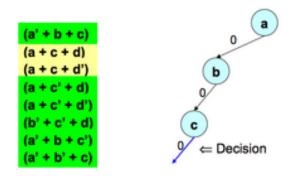


Figure: Basic DLL (Courtesy: Lintao Zhang)

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Backtracking Search (6/21)

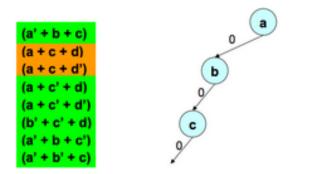


Figure: Basic DLL (Courtesy: Lintao Zhang)

Backtracking Search (7/21)

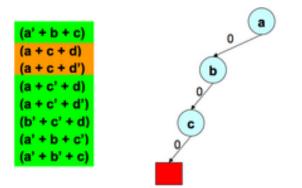


Figure: Basic DLL (Courtesy: Lintao Zhang)

Backtracking Search (8/21)

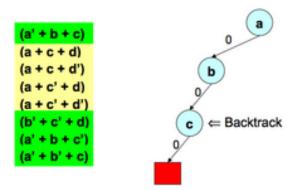


Figure: Basic DLL (Courtesy: Lintao Zhang)

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Backtracking Search (9/21)

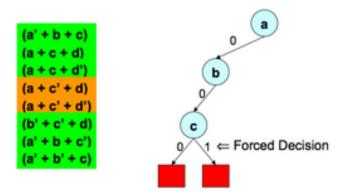


Figure: Basic DLL (Courtesy: Lintao Zhang)

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Backtracking Search (10/21)

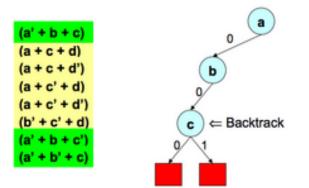


Figure: Basic DLL (Courtesy: Lintao Zhang)

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Backtracking Search (11/21)

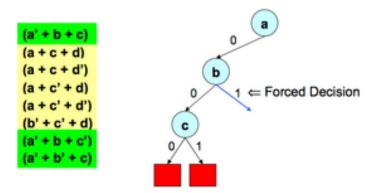


Figure: Basic DLL (Courtesy: Lintao Zhang)

Backtracking Search (12/21)

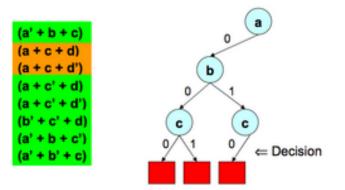


Figure: Basic DLL (Courtesy: Lintao Zhang)

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Backtracking Search (13/21)

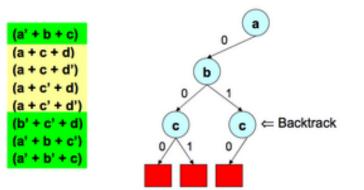


Figure: Basic DLL (Courtesy: Lintao Zhang)

Backtracking Search (14/21)

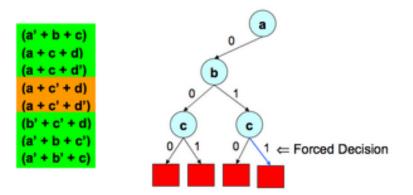
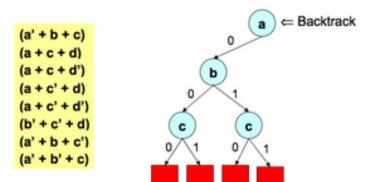


Figure: Basic DLL (Courtesy: Lintao Zhang)

Backtracking Search (15/21)



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Figure: Basic DLL (Courtesy: Lintao Zhang)

Backtracking Search (16/21)

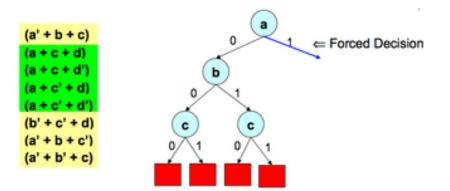


Figure: Basic DLL (Courtesy: Lintao Zhang)

Backtracking Search (17/21)

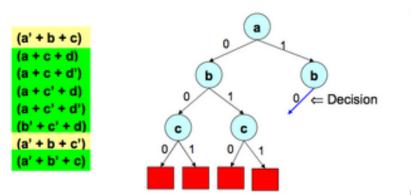
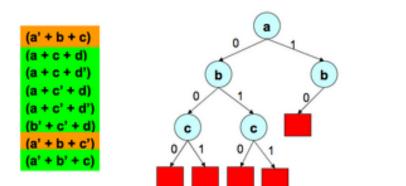


Figure: Basic DLL (Courtesy: Lintao Zhang)

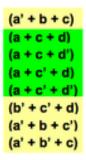
Backtracking Search (18/21)



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Figure: Basic DLL (Courtesy: Lintao Zhang)

Backtracking Search (19/21)



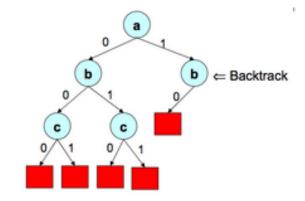


Figure: Basic DLL (Courtesy: Lintao Zhang)

Backtracking Search (20/21)

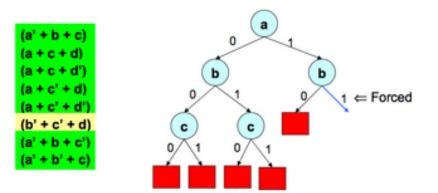


Figure: Basic DLL (Courtesy: Lintao Zhang)

Backtracking Search (21/21)

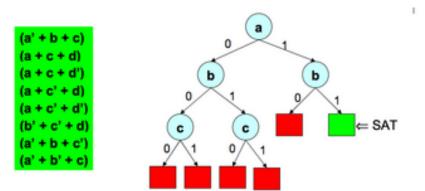


Figure: Basic DLL (Courtesy: Lintao Zhang)

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Backtracking Search

Don't build whole tree, but lazily search solutions

- Choose a variable x, set to True
- Remove constraints where x appears
- Recurse on remaining constraints
- Backtrack if a contradiction is found

(whew!)

DFS avoids space blowup (only need to save stack)

... but not time (natch)

Decision Procedures: Propositional Logic

Basics

- Propositional Logic 101
- Conjunctive Normal Form

Algorithms

- Resolution
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- Boolean Constraint Propagation
- Conflict Driven Learning & Backjumping

Often, we don't really have a choice...

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Unit Clause Rule

- ▶ If an (unsatisfied) Clause has one unassigned Literal
- Then that Literal must be True in any SAT assignment

Example

- Formula $(x_1 \lor \neg x_2 \lor x_3) \land (x_2 \lor \neg x_3) \land (\neg x_1 \lor \neg x_3)$
- Assignment $x_1 = T, x_2 = T$
- The last clause is a unit clause
- Any SAT assigment **must** set $\neg x_3 = T$ (i.e. $x_3 = F$)

Unit Clause Rule

- ▶ If an (unsatisfied) Clause has **one** unassigned Literal
- ▶ Then that Literal must be True in any SAT assignment

BCP or Unit Propagation

- Repeat applying unit clause rule
- Until no unit clause remains.

Boolean Constraint Propagation: Example

Revisit Example With BCP

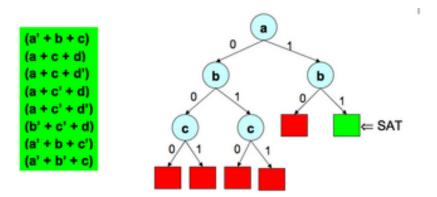


Figure: Boolean Constraint Propagation (Courtesy: Lintao Zhang)

DPLL = Backtracking Search + BCP

- Backtracking: Avoids space blowup
- BCP: Avoid doing obvious work
- Still repeatedly explore all choices (e.g. whole left subtree)

Wanted

- Means to *learn* to repeat *dead ends*
- Key to scaling to practical problems

Decision Procedures: Propositional Logic

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Conflict Driven Learning

Key Insight

- On finding conflict, don't (just) backtrack
- Learn new clause to prevent same conflict in future

Major breakthrough

- ► J. P. Marques-Silva and K. A. Sakallah, "GRASP A New Search Algorithm for Satisfiability," Proc. ICCAD 1996.
- R. J. Bayardo Jr. and R. C. Schrag "Using CSP look-back techniques to solve real world SAT instances." Proc. AAAI, 1997

Conflict Driven Learning

Resolve on conflict variable to learn new conflict clause

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- Add clause to set of clauses
- Backjump using conflict clause

Conflict Driven Learning Revisit Example With CDL

- Learn, Add, Backjump
- Vastly faster search

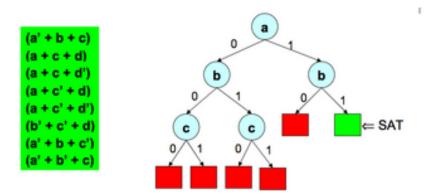


Figure: Boolean Constraint Propagation (Courtesy: Lintao Zhang)

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Backtracking Only (01/26)



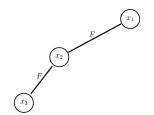
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Backtracking Only (02/26)

 x_1 F

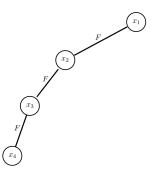
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Backtracking Only (03/26)



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Backtracking Only (04/26)

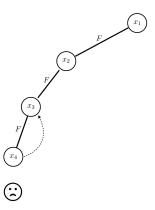


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Backtracking Only (05/26)

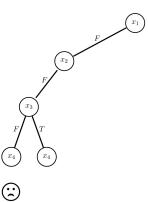
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|--------------|------------|---|------------|---|------------|-------------|
| | x_1 | V | x_3 | V | x_4 | F |
| | x_1 | V | x_3 | V | $\neg x_4$ | |
| \checkmark | x_1 | V | $\neg x_3$ | V | x_4 | |
| \checkmark | x_1 | V | $\neg x_3$ | V | $\neg x_4$ | F |
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| \checkmark | $\neg x_1$ | V | x_2 | V | $\neg x_3$ | |
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Backtracking Only (06/26)



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Backtracking Only (07/26)



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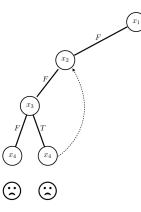
Backtracking Only (08/26)

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|--------------|------------|---|------------|---|------------|--------|-------------------------------------|---------------------------------------|---|---|---|---|---|---|---|---|
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| | x_1 | V | $\neg x_3$ | V | x_4 | | | | | | | | | | | |
| | x_1 | V | $\neg x_3$ | V | $\neg x_4$ | | | | F | F | F | F | F | F | F | F |
| \checkmark | $\neg x_2$ | V | $\neg x_3$ | V | x_4 | | | | | | | | | | | |
| \checkmark | $\neg x_1$ | V | x_2 | V | $\neg x_3$ | | G | (x_3) | (x_2) |
| \checkmark | $\neg x_1$ | V | $\neg x_2$ | V | x_3 | | | ý, | \bigvee | | | $\bigcup_{n=3}^{\infty}$ | | | $\bigcup_{n=3}^{\infty}$ | |
| | | | | | | | F | $_{\rm F}/$ | $_{r}/\lambda_{T}$ | $_{\rm F}/\lambda_{\rm T}$ | $_{r}/\lambda_{T}$ | $_{r}/ \setminus_{T}$ | $_{\rm F}/\lambda_{\rm T}$ | $_{r}/\lambda_{T}$ | $_{r}/ \setminus_{T}$ | $_{\rm F}/\lambda_{\rm T}$ |
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 x_1

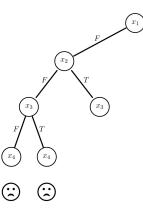
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Backtracking Only (09/26)



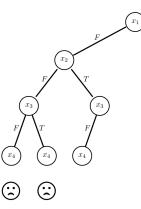
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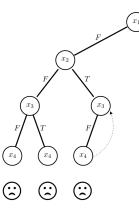
Backtracking Only (12/26)

| \checkmark | $\neg x_1$ | V | x_2 | V | x_3 | | | - / |
|--------------|------------|--------|------------|--------|------------|-------------------------------------|-------------------------------------|-------------------------------------|
| | x_1 | V | x_3 | V | x_4 |] | | F |
| | x_1 | \vee | x_3 | V | $\neg x_4$ | | (a | |
| \checkmark | x_1 | V | $\neg x_3$ | V | x_4 | | | 2 |
| \checkmark | x_1 | V | $\neg x_3$ | \vee | $\neg x_4$ | | F | \mathbf{n} |
| \checkmark | $\neg x_2$ | V | $\neg x_3$ | V | x_4 | | | \backslash |
| \checkmark | $\neg x_1$ | V | x_2 | V | $\neg x_3$ | G | r3 | (x_3) |
| \checkmark | $\neg x_1$ | V | $\neg x_2$ | V | x_3 | ζ | 2 | \bigcirc |
| | | | | | | F | Λ_T | $_{F}$ |
| | | | | | | ./ | ľ | 1 |
| | | | | | | \mathcal{A} | X | \mathcal{A} |
| | | | | | | $\begin{pmatrix} x_4 \end{pmatrix}$ | $\begin{pmatrix} x_4 \end{pmatrix}$ | $\begin{pmatrix} x_4 \end{pmatrix}$ |
| | | | | | | Ŭ | Ŭ | 0 |
| | | | | | | \bigcirc | \bigcirc | \bigcirc |
| | | | | | | G | C | S |
| | | | | | | \odot | \odot | \odot |

 x_1

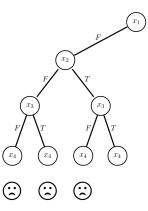
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Backtracking Only (13/26)



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Backtracking Only (14/26)



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Backtracking Only (15/26)

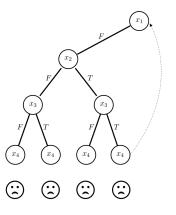
 $\neg x_1$ x_1 x_1 x_1 x_1 $\overline{\neg x_2}$

 $\neg x_1$

| $\lor x_2 \lor x_3$ | | | | |
|--|-------------------------------------|----------------|--------------|--------|
| $\lor x_3 \lor x_4$ | | | F | |
| $ \begin{array}{c c} \lor & x_3 & \lor & \neg x_4 \\ \hline \lor & \neg x_3 & \lor & x_4 \end{array} $ | | | | |
| $\bigvee \neg x_3 \lor \neg x_4$ $\lor \neg x_3 \lor x_4$ | | F | | |
| $\lor x_2 \lor \neg x_3$ | (x | 3 | \sum_{x} | 3 |
| $\lor \neg x_2 \lor x_3$ | 7 | イ | \mathbf{Y} | イ |
| | F | \backslash^T | F | T |
| | $\begin{pmatrix} x_4 \end{pmatrix}$ | x_4 | | |
| | \sim | \sim | \sim | \sim |
| | (::) | (::) | (::) | (::) |

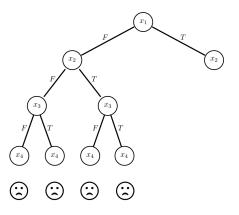
 x_1

Backtracking Only (16/26)

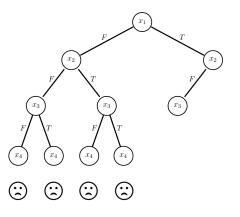


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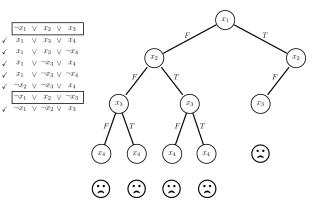
Backtracking Only (17/26)



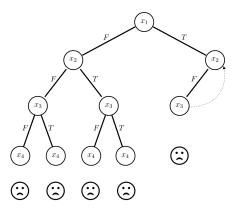
Backtracking Only (18/26)



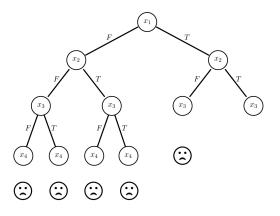
Backtracking Only (19/26)



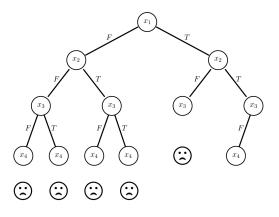
Backtracking Only (20/26)



Backtracking Only (21/26)



Backtracking Only (22/26)

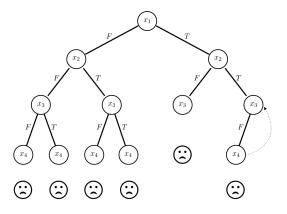


Backtracking Only (23/26)

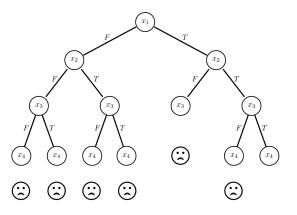
 x_1 $\neg x_1 \lor x_2 \lor x_3$ $\lor x_3 \lor x_4$ x_1 $x_1 \lor x_3 \lor \neg x_4$ x_2 x_2 $x_1 \lor \neg x_3 \lor x_4$ $x_1 \lor \neg x_3 \lor \neg x_4$ $\neg x_2 \lor \neg x_3 \lor x_4$ $\neg x_1 \lor x_2 \lor \neg x_3$ x_3 x_3 x_3 x_3 $\neg x_1 \lor \neg x_2 \lor x_3$ \mathbb{C} x_4 x_4 x_4 x_4 x_4 \odot (\dot{z}) (:)(:)....

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Backtracking Only (24/26)

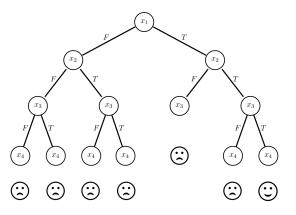


Backtracking Only (25/26)



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Backtracking Only (26/26)



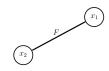
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Boolean Constraint Propagation (01/23)



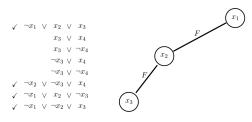
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Boolean Constraint Propagation (02/23)

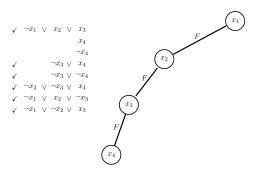


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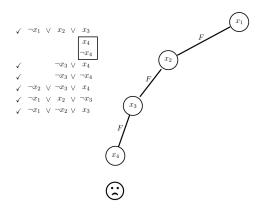
Boolean Constraint Propagation (03/23)



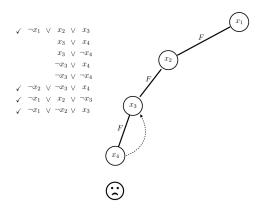
Boolean Constraint Propagation (04/23)



Boolean Constraint Propagation (05/23)

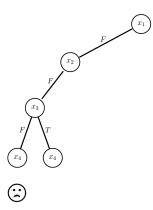


Boolean Constraint Propagation (06/23)



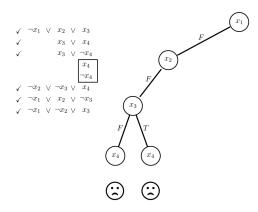
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Boolean Constraint Propagation (07/23)

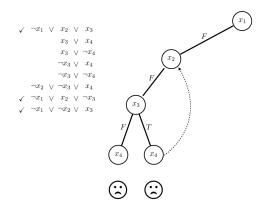


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Boolean Constraint Propagation (08/23)

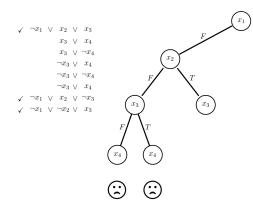


Boolean Constraint Propagation (09/23)



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Boolean Constraint Propagation (10/23)



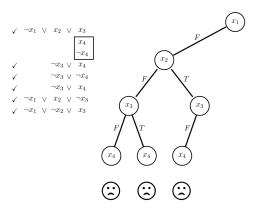
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Boolean Constraint Propagation (11/23)

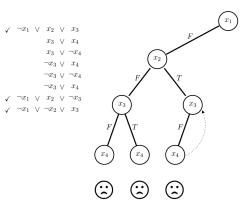
 $\neg x_1 \lor x_2 \lor x_3$ x_4 $\neg x_A$ $\neg x_3 \lor x_4$ $\neg x_3 \lor \neg x_4$ $\neg x_3 \lor x_4$ $\lor x_2 \lor \neg x_3$ x_3 x_3 $\neg x_1 \lor \neg x_2 \lor x_3$ x_4 x_4 x_4 (\dot{z}) (\dot{z})

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Boolean Constraint Propagation (12/23)

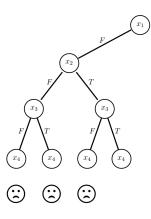


Boolean Constraint Propagation (13/23)

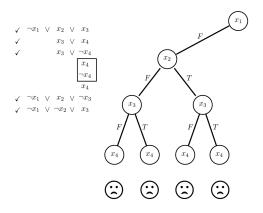


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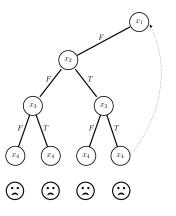
Boolean Constraint Propagation (14/23)



Boolean Constraint Propagation (15/23)

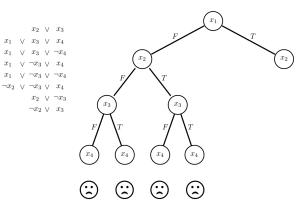


Boolean Constraint Propagation (16/23)



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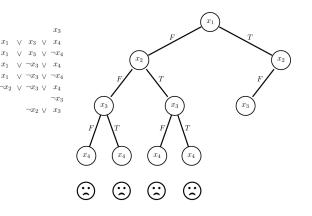
Boolean Constraint Propagation (17/23)



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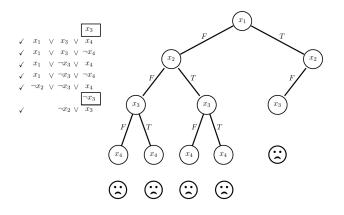
Boolean Constraint Propagation (18/23)

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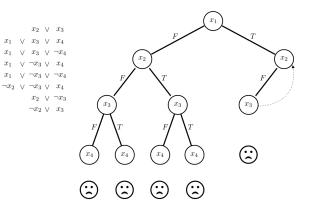


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Boolean Constraint Propagation (19/23)

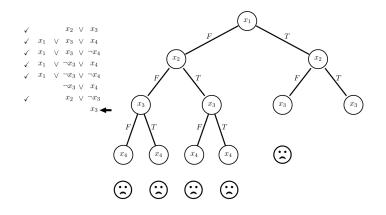


Boolean Constraint Propagation (20/23)



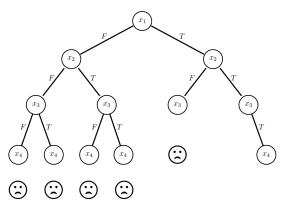
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Boolean Constraint Propagation (21/23)



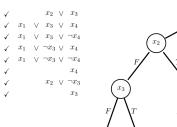
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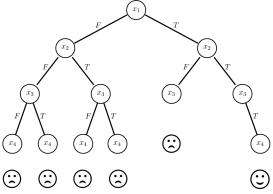
Boolean Constraint Propagation (22/23)



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Boolean Constraint Propagation (23/23)





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Conflict Driven Learning (01/21)



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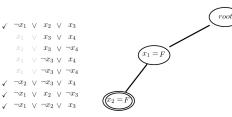
Conflict Driven Learning (02/21)

root

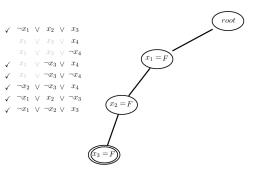
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 $\checkmark \neg x_1 \lor \neg x_2 \lor x_3$

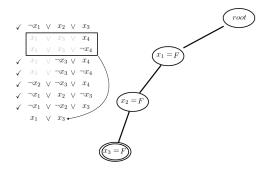
Conflict Driven Learning (03/21)



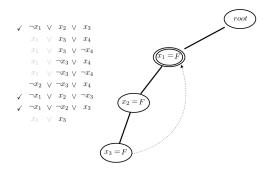
Conflict Driven Learning (04/21)



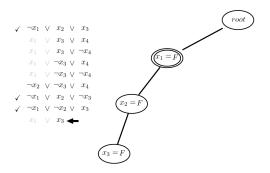
Conflict Driven Learning (05/21)



Conflict Driven Learning (06/21)

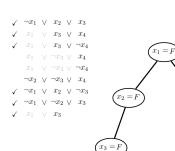


Conflict Driven Learning (07/21)

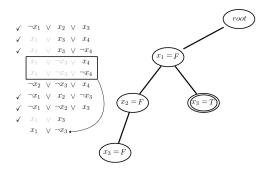


Conflict Driven Learning (08/21)

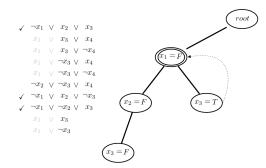
root



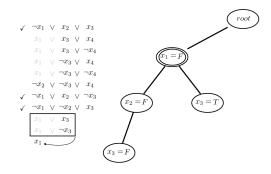
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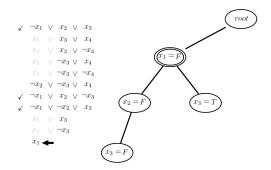
Conflict Driven Learning (10/21)



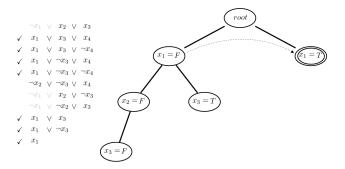
Conflict Driven Learning (11/21)



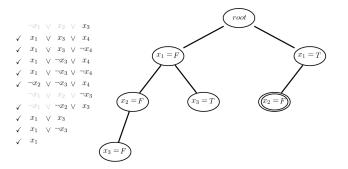
Conflict Driven Learning (12/21)



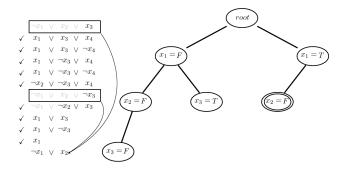
Conflict Driven Learning (13/21)



Conflict Driven Learning (14/21)

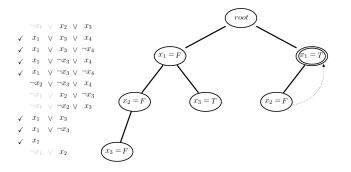


Conflict Driven Learning (15/21)



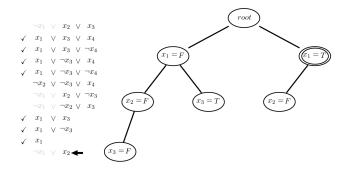
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Conflict Driven Learning (16/21)



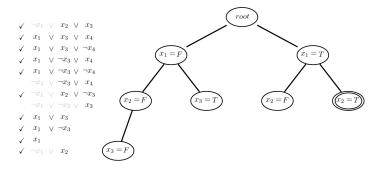
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Conflict Driven Learning (17/21)

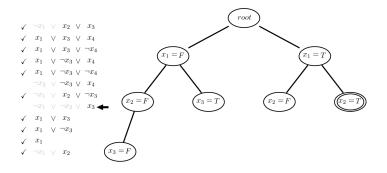


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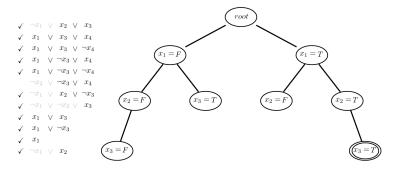
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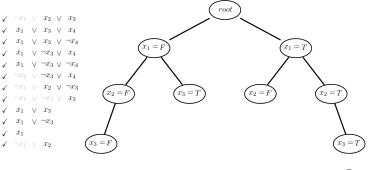
Conflict Driven Learning (19/21)



Conflict Driven Learning (20/21)



Conflict Driven Learning (21/21)



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More Details about SAT Solvers

Lectures By Lintao Zhang (ZChaff)

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Next Time: SMT = SAT + Theories

- 1. Propositional Logic
- 2. Combining Theories
 - Equality + Uninterpreted Functions

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- Difference-Bounded Arithmetic
- 3. Combining SAT + Theories