## SAT Solvers

Ranjit Jhala, UC San Diego

April 9, 2013

## Decision Procedures

We will look very closely at the following

1. Propositional Logic
2. Theory of Equality
3. Theory of Uninterpreted Functions
4. Theory of Difference-Bounded Arithmetic

Decision Problem: Satisfaction

- Does eval s p return True for some assignment s?
- "Can we assign the variables to make the formula true" ?


## Decision Procedures

We will look very closely at the following

1. Propositional Logic
2. Theory of Equality
3. Theory of Uninterpreted Functions
4. Theory of Difference-Bounded Arithmetic

Why?

- Representative
- Have "efficient" algorithms


## Decision Procedures

We will look very closely at the following

1. Propositional Logic
2. Theory of Equality
3. Theory of Uninterpreted Functions
4. Theory of Difference-Bounded Arithmetic

Plan

- First in isolation
- Then in combination
- Very slick SW-Eng, based on logic


## Decision Procedures: Propositional Logic

Popularly called SAT Solvers

## Decision Procedures: Propositional Logic

Basics

- Propositional Logic 101
- Conjunctive Normal Form
- Resolution

Algorithms

- Resolution
- Backtracking Search
- Boolean Constraint Propagation
- Conflict Driven Learning \& Backjumping


## Decision Procedures: Propositional Logic

Basics

- Propositional Logic 101
- Conjunctive Normal Form
- Resolution

Algorithms

- Resolution
- Backtracking Search
- Boolean Constraint Propagation
- Conflict Driven Learning \& Backjumping


## Propositional Logic 101

## Propositional Variables

data PVar

Propositional Formulas
data Formula $=$ Prop PVar
Not Formula
Formula 'And' Formula
Formula 'Or' Formula

## Decision Procedures: Propositional Logic

Basics

- Propositional Logic 101
- Conjunctive Normal Form
- Resolution

Algorithms

- Resolution
- Backtracking Search
- Boolean Constraint Propagation
- Conflict Driven Learning \& Backjumping


## Conjunctive Normal Form

Restricted representation of Formula
Literals: Variables or Negated Variables
data Literal = Pos PVar | Neg PVar

Clauses: Disjunctions (Or) of Literals
data Clauses = [Literal]

CNF Formulas: Conjunctions (And) of Clauses
data CnfFormula $=$ [Clauses]

## Conjunctive Normal Form: Example

Consider a Formula
$\left(x_{1} \vee x_{2}\right) \wedge\left(\neg x_{1} \vee x_{3}\right) \wedge \neg x_{3}$
Represented as a Formula

|  | (Prop 1 |
| :--- | :--- |
| 'And' (Not (Prop 1) | 'Or' Prop 2) |
| 'And' (Not (Prop 3) |  |

Represented as a CnfFormula
[ [Pos 1 , Pos 2]
, $[\operatorname{Neg} 1$, Pos 3]
, $[\operatorname{Neg} 3]]$

## Conjunctive Normal Form Conversion

Theorem There is a poly-time function
toCNF :: Formula -> CnfFormula
toCNF = error "Exercise For The Reader"
Such that any f is satisfiable iff ( toCNF f ) is satisfiable.

- toCNF adds new variables for sub-formulas
- otherwise, an exponential blowup in CnfFormula size


## Conjunctive Normal Form Conversion

Theorem There is a poly-time function
toCNF :: Formula -> CnfFormula
toCNF = error "Exercise For The Reader"
Such that any f is satisfiable iff ( toCNF f ) is satisfiable.
Henceforth Only consider formulas in Conjunctive Normal Form
Formulas

## Decision Procedures: Propositional Logic

Basics

- Propositional Logic 101
- Conjunctive Normal Form

Algorithms

- Resolution
- Backtracking Search
- Boolean Constraint Propagation
- Conflict Driven Learning \& Backjumping


## Properties of CNF

## Pure Variable

- One which appears only $+v e$ or $-v e$ in a CnfFormula


## Empty Clause

- If a CnfFormula has some Clause without Literals
- Then the CnfFormula is UNSAT

Trivial Formula

- If a CnfFormula has no Clause
- Or every variable is pure
- Then the CnfFormula is SAT


## Goal

Determine satisfaction by reducing CnfFormula to one of

- Empty Clause (ie UNSAT), or
- Trivial Formula (ie SAT).


## Reducing Formulas By Resolution

("Reduce" is, perhaps, not the best word...)
Resolution: For any $A, B$ and variable $x$, the formula

$$
(A \vee x) \wedge(B \vee \neg x)
$$

is equivalent to the formula

$$
(A \vee B)
$$

- The variable $x$ is called a pivot variable


## General Resolution

Resolution: For any $A_{i}, B_{j}$ and variable $x$, the formula

$$
\bigwedge_{i}\left(A_{i} \vee x\right) \wedge \bigwedge_{j}\left(B_{j} \vee \neg x\right)
$$

is equivalent to the formula

$$
\bigwedge_{i, j}\left(A_{i} \vee B_{j}\right)
$$

- Pivot variable $x$ is eliminated by resolution


## Davis-Putnam Algorithm: Example 1

Input Formula

- $\left.\left(x_{1} \vee x_{2} \vee x_{3}\right) \wedge\left(x_{2} \vee \neg x_{3} \vee x_{5}\right) \wedge\left(\neg x_{2} \vee x_{4}\right)\right)$

Pivot on $x_{2}$

- $\left(x_{1} \vee x_{3} \vee x_{4}\right) \wedge\left(\neg x_{3} \vee x_{5} \vee x_{4}\right)$

Pivot on $x_{3}$

- $\left(x_{1} \vee x_{4} \vee x_{5}\right)$

All variables are pure ... hence, SAT

## Davis-Putnam Algorithm: Example 2

Input Formula

- $\left(x_{1} \vee x_{2}\right) \wedge\left(x_{1} \vee \neg x_{2}\right) \wedge\left(\neg x_{1} \vee x_{3}\right) \wedge\left(\neg x_{1} \vee \neg x_{3}\right)$

Pivot on $x_{2}$

- $\left(x_{1}\right) \wedge\left(\neg x_{1} \vee x_{3}\right) \wedge\left(\neg x_{1} \vee \neg x_{3}\right)$

Pivot on $x_{3}$

- $\left(x_{1}\right) \wedge\left(\neg x_{1}\right)$

Pivot on $x_{1}$

- ()

Empty clause . . . hence, UNSAT

## Davis-Putnam Algorithm

Algorithm

1. Select pivot and perform resolution
2. Repeat until SAT or UNSAT

## Issues?

- Space blowup (formula size blows up on resolution)


## Decision Procedures: Propositional Logic

Basics

- Propositional Logic 101
- Conjunctive Normal Form

Algorithms

- Resolution
- Backtracking Search
- Boolean Constraint Propagation
- Conflict Driven Learning \& Backjumping


## Decision Tree: Describes Space of All Assignments



Figure: SAT Decision Tree (Courtesy: Lintao Zhang)

Decision Tree: SAT via Depth First Search


Figure: DFS On Decision Tree (Courtesy: Lintao Zhang)

## Backtracking Search

Don't build whole tree, but lazily search solutions

- Choose a variable $x$, set to True
- Remove constraints where $x$ appears
- Recurse on remaining constraints
- Backtrack if a contradiction is found


## Backtracking Search (1/21)

( $a^{\prime}+b+c$ )<br>$(a+c+d)$<br>$\left(a+c+d^{\prime}\right)$<br>$\left(a+c^{\prime}+d\right)$<br>$\left(a+c^{\prime}+d^{\prime}\right)$<br>(b' $+c^{\prime}+d$ )<br>( $\mathbf{a}^{\prime}+\mathrm{b}+\mathrm{c}^{\prime}$ )<br>( $a^{\prime}+b^{\prime}+c$ )

Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (2/21)

```
(a'+b+c)
(a+c+d)
(a+c+d')
(a+c'+d)
(a+c' + d')
(b' + c' + d)
(a' + b + c')
(a'+b'+c)
```

Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (3/21)

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+c^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$

Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (4/21)

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+c^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$



Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (5/21)

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+c^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$



Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (6/21)

$\left(a^{\prime}+b+c\right)$
$(a+c+d)$
$\left(a+c+d^{\prime}\right)$
$\left(a+c^{\prime}+d\right)$
$\left(a+c^{\prime}+d^{\prime}\right)$
$\left(b^{\prime}+c^{\prime}+d\right)$
$\left(a^{\prime}+b+c^{\prime}\right)$
$\left(a^{\prime}+b^{\prime}+c\right)$


Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (7/21)

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+c^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$



Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (8/21)

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+b^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$



Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (9/21)

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+c^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$



Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (10/21)

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+c^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$



Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (11/21)

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+c^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$



Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (12/21)

```
(a'+b+c)
(a+c+d)
(a+c+d')
(a+c'+d)
(a+c'+ d')
(b'+c'+d)
(a'+b+c')
(a'+b' + c)
```



Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (13/21)

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+b^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$



Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (14/21)

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+b^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$



Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (15/21)

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+c^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$



Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (16/21)

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+c^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$



Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (17/21)

$\left(a^{\prime}+b+c\right)$
$(a+c+d)$
$\left(a+c+d^{\prime}\right)$
$\left(a+c^{\prime}+d\right)$
$\left(a+c^{\prime}+d^{\prime}\right)$
$\left(b^{\prime}+c^{\prime}+d\right)$
$\left(a^{\prime}+b+c^{\prime}\right)$
$\left(a^{\prime}+b^{\prime}+c\right)$


Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (18/21)

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+c^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$



Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (19/21)

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+b^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$



Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (20/21)

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+c^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$



Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search (21/21)

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+b^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$



Figure: Basic DLL (Courtesy: Lintao Zhang)

## Backtracking Search

Don't build whole tree, but lazily search solutions

- Choose a variable $x$, set to True
- Remove constraints where $x$ appears
- Recurse on remaining constraints
- Backtrack if a contradiction is found
(whew!)
- DFS avoids space blowup (only need to save stack) ...
- ... but not time (natch)


## Decision Procedures: Propositional Logic

Basics

- Propositional Logic 101
- Conjunctive Normal Form

Algorithms

- Resolution
- Backtracking Search
- Boolean Constraint Propagation
- Conflict Driven Learning \& Backjumping


## Boolean Constraint Propagation

Often, we don't really have a choice...

## Boolean Constraint Propagation

## Unit Clause Rule

- If an (unsatisfied) Clause has one unassigned Literal
- Then that Literal must be True in any SAT assignment

Example

- Formula $\left(x_{1} \vee \neg x_{2} \vee x_{3}\right) \wedge\left(x_{2} \vee \neg x_{3}\right) \wedge\left(\neg x_{1} \vee \neg x_{3}\right)$
- Assignment $x_{1}=T, x_{2}=T$
- The last clause is a unit clause
- Any SAT assigment must set $\neg x_{3}=T$ (i.e. $x_{3}=F$ )


## Boolean Constraint Propagation

## Unit Clause Rule

- If an (unsatisfied) Clause has one unassigned Literal
- Then that Literal must be True in any SAT assignment

BCP or Unit Propagation

- Repeat applying unit clause rule
- Until no unit clause remains.


## Boolean Constraint Propagation: Example

Revisit Example With BCP

```
(a'+b + c)
(a+c+d)
(a+c+d')
(a+c' +d)
(a+c'+ +')
(b' + c' + d)
(a'+b+c')
(a'+b'+c)
```



Figure: Boolean Constraint Propagation (Courtesy: Lintao Zhang)

## Boolean Constraint Propagation

DPLL $=$ Backtracking Search + BCP

- Backtracking: Avoids space blowup
- BCP: Avoid doing obvious work
- Still repeatedly explore all choices (e.g. whole left subtree)


## Wanted

- Means to learn to repeat dead ends
- Key to scaling to practical problems


## Decision Procedures: Propositional Logic

Basics

- Propositional Logic 101
- Conjunctive Normal Form

Algorithms

- Resolution
- Backtracking Search
- Boolean Constraint Propagation
- Conflict Driven Learning \& Backjumping


## Conflict Driven Learning

Key Insight

- On finding conflict, don't (just) backtrack
- Learn new clause to prevent same conflict in future

Major breakthrough

- J. P. Marques-Silva and K. A. Sakallah, "GRASP - A New Search Algorithm for Satisfiability," Proc. ICCAD 1996.
- R. J. Bayardo Jr. and R. C. Schrag "Using CSP look-back techniques to solve real world SAT instances." Proc. AAAI, 1997


## Conflict Driven Learning

- Resolve on conflict variable to learn new conflict clause
- Add clause to set of clauses
- Backjump using conflict clause


## Conflict Driven Learning

Revisit Example With CDL

- Learn, Add, Backjump
- Vastly faster search

$$
\begin{aligned}
& \left(a^{\prime}+b+c\right) \\
& (a+c+d) \\
& \left(a+c+d^{\prime}\right) \\
& \left(a+c^{\prime}+d\right) \\
& \left(a+c^{\prime}+d^{\prime}\right) \\
& \left(b^{\prime}+c^{\prime}+d\right) \\
& \left(a^{\prime}+b+b^{\prime}\right) \\
& \left(a^{\prime}+b^{\prime}+c\right)
\end{aligned}
$$



Figure: Boolean Constraint Propagation (Courtesy: Lintao Zhang)

## Backtracking Only (01/26)

$$
\begin{array}{ccccc}
\neg x_{1} & \vee & x_{2} & \vee & x_{3}  \tag{1}\\
x_{1} & \vee & x_{3} & \vee & x_{4} \\
x_{1} & \vee & x_{3} & \vee & \neg x_{4} \\
x_{1} & \vee & \neg x_{3} & \vee & x_{4} \\
x_{1} & \vee & \neg x_{3} & \vee & \neg x_{4} \\
\neg x_{2} & \vee & \neg x_{3} & \vee & x_{4} \\
\neg x_{1} & \vee & x_{2} & \vee & \neg x_{3} \\
\neg x_{1} & \vee & \neg x_{2} & \vee & x_{3}
\end{array}
$$

## Backtracking Only $(02 / 26)$

| $\checkmark$ | $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $x_{3}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | $x_{1}$ | $\vee$ | $x_{3}$ | $\vee$ | $x_{4}$ |
| $x_{1}$ | $\vee$ | $x_{3}$ | $\vee$ | $\neg x_{4}$ |  |
| $x_{1}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $x_{4}$ |  |
|  | $x_{1}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $\neg x_{4}$ |
|  | $\neg x_{2}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $x_{4}$ |
| $\checkmark$ | $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $\neg x_{3}$ |
| $\checkmark$ | $\neg x_{1}$ | $\vee$ | $\neg x_{2}$ | $\vee$ | $x_{3}$ |



## Backtracking Only (03/26)

$$
\begin{array}{cccccc}
\checkmark & \neg x_{1} & \vee & x_{2} & \vee & x_{3} \\
x_{1} & \vee & x_{3} & \vee & x_{4} \\
x_{1} & \vee & x_{3} & \vee & \neg x_{4} \\
x_{1} & \vee & \neg x_{3} & \vee & x_{4} \\
x_{1} & \vee & \neg x_{3} & \vee & \neg x_{4} \\
\checkmark & \neg x_{2} & \vee & \neg x_{3} & \vee & x_{4} \\
\checkmark & \neg x_{1} & \vee & x_{2} & \vee & \neg x_{3} \\
\checkmark & \neg x_{1} & \vee & \neg x_{2} & \vee & x_{3}
\end{array}
$$



## Backtracking Only (04/26)



## Backtracking Only (05/26)


©

## Backtracking Only (06/26)


©

## Backtracking Only (07/26)


©

## Backtracking Only (08/26)



## Backtracking Only (09/26)

$\checkmark$| $\checkmark$ | $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $x_{3}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | $x_{1}$ | $\vee$ | $x_{3}$ | $\vee$ | $x_{4}$ |
|  | $x_{1}$ | $\vee$ | $x_{3}$ | $\vee$ | $\neg x_{4}$ |
|  | $x_{1}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $x_{4}$ |
|  | $x_{1}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $\neg x_{4}$ |
|  | $\neg x_{2}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $x_{4}$ |
| $\checkmark$ | $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $\neg x_{3}$ |
| $\checkmark$ | $\neg x_{1}$ | $\vee$ | $\neg x_{2}$ | $\vee$ | $x_{3}$ |


; ;

## Backtracking Only (10/26)

$$
\checkmark \begin{array}{ccccc}
\neg x_{1} & \vee & x_{2} & \vee & x_{3} \\
x_{1} & \vee & x_{3} & \vee & x_{4} \\
x_{1} & \vee & x_{3} & \vee & \neg x_{4} \\
x_{1} & \vee & \neg x_{3} & \vee & x_{4} \\
x_{1} & \vee & \neg x_{3} & \vee & \neg x_{4} \\
& \neg x_{2} & \vee & \neg x_{3} & \vee \\
\checkmark & x_{4} \\
\checkmark & \neg x_{1} & \vee & x_{2} & \vee \\
\checkmark & \neg x_{3} \\
\checkmark & \neg x_{1} & \vee & \neg x_{2} & \vee \\
& x_{3}
\end{array}
$$


$\infty$

## Backtracking Only (11/26)

$$
\begin{array}{cccccc}
\checkmark & \neg x_{1} & \vee & x_{2} & \vee & x_{3} \\
& x_{1} & \vee & x_{3} & \vee & x_{4} \\
& x_{1} & \vee & x_{3} & \vee & \neg x_{4} \\
\checkmark & x_{1} & \vee & \neg x_{3} & \vee & x_{4} \\
\checkmark & x_{1} & \vee & \neg x_{3} & \vee & \neg x_{4} \\
\checkmark & \neg x_{2} & \vee & \neg x_{3} & \vee & x_{4} \\
\checkmark & \neg x_{1} & \vee & x_{2} & \vee & \neg x_{3} \\
\checkmark & \neg x_{1} & \vee & \neg x_{2} & \vee & x_{3}
\end{array}
$$


$\infty$

## Backtracking Only (12/26)


$\infty$ -

## Backtracking Only (13/26)

| $\checkmark$ | $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $x_{3}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | $x_{1}$ | $\vee$ | $x_{3}$ | $\vee$ | $x_{4}$ |
|  | $x_{1}$ | $\vee$ | $x_{3}$ | $\vee$ | $\neg x_{4}$ |
|  | $x_{1}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $x_{4}$ |
|  | $x_{1}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $\neg x_{4}$ |
|  | $\neg x_{2}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $x_{4}$ |
| $\checkmark$ | $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $\neg x_{3}$ |
| $\checkmark$ | $\neg x_{1}$ | $\vee$ | $\neg x_{2}$ | $\vee$ | $x_{3}$ |


$\infty$ -

## Backtracking Only (14/26)

|  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\checkmark$ | $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $x_{3}$ |
| $\checkmark$ | $x_{1}$ | $\vee$ | $x_{3}$ | $\vee$ | $x_{4}$ |
| $\checkmark$ | $x_{1}$ | $\vee$ | $x_{3}$ | $\vee$ | $\neg x_{4}$ |
|  | $x_{1}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $x_{4}$ |
|  | $x_{1}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $\neg x_{4}$ |
|  | $\neg x_{2}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $x_{4}$ |
| $\checkmark$ | $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $\neg x_{3}$ |
| $\checkmark$ | $\neg x_{1}$ | $\vee$ | $\neg x_{2}$ | $\vee$ | $x_{3}$ |


$\infty$ -

## Backtracking Only (15/26)



## Backtracking Only (16/26)

|  |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- |
| $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $x_{3}$ |
| $x_{1}$ | $\vee$ | $x_{3}$ | $\vee$ | $x_{4}$ |
| $x_{1}$ | $\vee$ | $x_{3}$ | $\vee$ | $\neg x_{4}$ |
| $x_{1}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $x_{4}$ |
| $x_{1}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $\neg x_{4}$ |
| $\neg x_{2}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $x_{4}$ |
| $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $\neg x_{3}$ |
| $\neg x_{1}$ | $\vee$ | $\neg x_{2}$ | $\vee$ | $x_{3}$ |


$\cdots$ -

Backtracking Only (17/26)


## Backtracking Only (18/26)

|  | $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $x_{3}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\checkmark$ | $x_{1}$ | $\vee$ | $x_{3}$ | $\vee$ | $x_{4}$ |
| $\checkmark$ | $x_{1}$ | $\vee$ | $x_{3}$ | $\vee$ | $\neg x_{4}$ |
| $\checkmark$ | $x_{1}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $x_{4}$ |
| $\checkmark$ | $x_{1}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $\neg x_{4}$ |
| $\checkmark$ | $\neg x_{2}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $x_{4}$ |
|  | $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $\neg x_{3}$ |
| $\checkmark$ | $\neg x_{1}$ | $\vee$ | $\neg x_{2}$ | $\vee$ | $x_{3}$ |


$\cdots$ -

## Backtracking Only (19/26)



## Backtracking Only (20/26)



Backtracking Only (21/26)


Backtracking Only (22/26)


Backtracking Only (23/26)


Backtracking Only (24/26)


Backtracking Only (25/26)


Backtracking Only (26/26)


## Boolean Constraint Propagation $(01 / 23)$

| $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $x_{3}$ |
| :---: | :---: | :---: | :---: | :---: |
| $x_{1}$ | $\vee$ | $x_{3}$ | $\vee$ | $x_{4}$ |
| $x_{1}$ | $\vee$ | $x_{3}$ | $\vee$ | $\neg x_{4}$ |
| $x_{1}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $x_{4}$ |
| $x_{1}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $\neg x_{4}$ |
| $\neg x_{2}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $x_{4}$ |
| $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $\neg x_{3}$ |
| $\neg x_{1}$ | $\vee$ | $\neg x_{2}$ | $\vee$ | $x_{3}$ |

Boolean Constraint Propagation $(02 / 23)$

| $\checkmark$ | $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $x_{3}$ |
| :--- | :--- | :--- | :--- | :--- | :--- |
|  |  | $x_{3}$ | $\vee$ | $x_{4}$ |  |
|  |  | $x_{3}$ | $\vee$ | $\neg x_{4}$ |  |
|  |  |  | $\neg x_{3}$ | $\vee$ | $x_{4}$ |
|  |  | $\neg x_{3}$ | $\vee$ | $\neg x_{4}$ |  |
|  |  | $\neg x_{2}$ | $\vee$ | $\neg x_{3}$ | $\vee$ |$x_{4}$



## Boolean Constraint Propagation $(03 / 23)$



## Boolean Constraint Propagation $(04 / 23)$



## Boolean Constraint Propagation $(05 / 23)$


©

## Boolean Constraint Propagation $(06 / 23)$


©

## Boolean Constraint Propagation $(07 / 23)$


©

## Boolean Constraint Propagation $(08 / 23)$


; ;

## Boolean Constraint Propagation $(09 / 23)$


; ;

## Boolean Constraint Propagation (10/23)


; $;$

## Boolean Constraint Propagation $(11 / 23)$


; $;$

## Boolean Constraint Propagation $(12 / 23)$



## Boolean Constraint Propagation $(13 / 23)$

$$
\begin{array}{llllll}
\checkmark & \neg x_{1} & \vee & x_{2} & \vee & x_{3} \\
& & x_{3} & \vee & x_{4} \\
& & x_{3} & \vee & \neg x_{4} \\
& & \neg x_{3} & \vee & x_{4} \\
& & \neg x_{3} & \vee & \neg x_{4} \\
& & \neg x_{3} & \vee & x_{4} \\
& & & & & \\
& & x_{1} & \vee & & x_{2} \\
\checkmark & \neg x_{3} \\
\checkmark & \neg x_{1} & \vee & \neg x_{2} & \vee & x_{3}
\end{array}
$$


; ; ;

## Boolean Constraint Propagation $(14 / 23)$



## Boolean Constraint Propagation $(15 / 23)$



## Boolean Constraint Propagation $(16 / 23)$

$$
\begin{array}{ccccc}
\neg x_{1} & \vee & x_{2} & \vee & x_{3} \\
x_{1} & \vee & x_{3} & \vee & x_{4} \\
x_{1} & \vee & x_{3} & \vee & \neg x_{4} \\
x_{1} & \vee & \neg x_{3} & \vee & x_{4} \\
x_{1} & \vee & \neg x_{3} & \vee & \neg x_{4} \\
\neg x_{2} & \vee & \neg x_{3} & \vee & x_{4} \\
\neg x_{1} & \vee & x_{2} & \vee & \neg x_{3} \\
\neg x_{1} & \vee & \neg x_{2} & \vee & x_{3}
\end{array}
$$


$\cdots$ -

## Boolean Constraint Propagation $(17 / 23)$



## Boolean Constraint Propagation $(18 / 23)$



## Boolean Constraint Propagation $(19 / 23)$



Boolean Constraint Propagation $(20 / 23)$


Boolean Constraint Propagation $(21 / 23)$


Boolean Constraint Propagation $(22 / 23)$


Boolean Constraint Propagation $(23 / 23)$


## Conflict Driven Learning (01/21)

| $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $x_{3}$ |
| :---: | :---: | :---: | :---: | :---: |
| $x_{1}$ | $\vee$ | $x_{3}$ | $\vee$ | $x_{4}$ |
| $x_{1}$ | $\vee$ | $x_{3}$ | $\vee$ | $\neg x_{4}$ |
| $x_{1}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $x_{4}$ |
| $x_{1}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $\neg x_{4}$ |
| $\neg x_{2}$ | $\vee$ | $\neg x_{3}$ | $\vee$ | $x_{4}$ |
| $\neg x_{1}$ | $\vee$ | $x_{2}$ | $\vee$ | $\neg x_{3}$ |
| $\neg x_{1}$ | $\vee$ | $\neg x_{2}$ | $\vee$ | $x_{3}$ |



## Conflict Driven Learning (02/21)

$$
\begin{array}{ccccc}
\checkmark x_{1} & \vee & x_{2} & \vee & x_{3} \\
x_{1} & \vee & x_{3} & \vee & x_{4} \\
x_{1} & \vee & x_{3} & \vee & \neg x_{4} \\
x_{1} & \vee & \neg x_{3} & \vee & x_{4} \\
x_{1} & \vee & \neg x_{3} & \vee & \neg x_{4} \\
\neg x_{2} & \vee & \neg x_{3} & \vee & x_{4} \\
\checkmark & \neg x_{1} & \vee & x_{2} & \vee \\
\checkmark & \neg x_{3} \\
\checkmark & \neg x_{1} & \vee & \neg x_{2} & \vee \\
& x_{3}
\end{array}
$$



## Conflict Driven Learning (03/21)



## Conflict Driven Learning (04/21)



## Conflict Driven Learning (05/21)



## Conflict Driven Learning (06/21)



## Conflict Driven Learning (07/21)



## Conflict Driven Learning (08/21)



## Conflict Driven Learning (09/21)



## Conflict Driven Learning (10/21)



## Conflict Driven Learning (11/21)



## Conflict Driven Learning (12/21)



## Conflict Driven Learning (13/21)



## Conflict Driven Learning (14/21)



## Conflict Driven Learning (15/21)



## Conflict Driven Learning (16/21)



## Conflict Driven Learning (17/21)



Conflict Driven Learning (18/21)


Conflict Driven Learning (19/21)


## Conflict Driven Learning (20/21)



## Conflict Driven Learning (21/21)



## More Details about SAT Solvers

Lectures By Lintao Zhang (ZChaff)

- 1
- 2


## Next Time: SMT = SAT + Theories

1. Propositional Logic
2. Combining Theories

- Equality + Uninterpreted Functions
- Difference-Bounded Arithmetic

3. Combining SAT + Theories
